### Certified CNF Translations for Pseudo-Boolean Solving

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SAT '22 paper joint with Stephan Gocht, Ruben Martins, and Andy Oertel

# Pseudo-Boolean (PB) Solving



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![](_page_4_Figure_1.jpeg)

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SAT solver

![](_page_6_Figure_1.jpeg)

![](_page_7_Figure_1.jpeg)

![](_page_8_Figure_1.jpeg)

![](_page_9_Figure_1.jpeg)

# Certifying Pseudo-Boolean Solver Results with Proof Logging?

![](_page_10_Figure_1.jpeg)

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![](_page_11_Figure_1.jpeg)

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# Certifying Pseudo-Boolean Solver Results with Proof Logging?

![](_page_12_Figure_1.jpeg)

- Correctness of SAT solver result can be certified [HHW13a, HHW13b, WHH14]
- PB-to-CNF translation uncertified!

 $\ell_1 + \ell_2 + \ell_3 \ge 2$ 

### $\ell_1+\ell_2+\ell_3\geq 2$

$\overline{\ell}_1 \vee s_{1,1}$	$\overline{\ell}_2 \vee \overline{s}_{1,1} \vee s_{2,2}$	$\ell_3 \vee s_{2,1} \vee \overline{s}_{3,1}$	$\overline{\ell}_3 \vee \overline{s}_{2,2} \vee s_{3,3}$	
$\ell_1 \vee \overline{s}_{1,1}$	$\ell_2 \vee \overline{s}_{2,2}$	$\overline{\ell}_3 \vee \overline{s}_{2,1} \vee s_{3,2}$	$\ell_3 \vee \overline{s}_{3,3}$	
$\overline{\ell}_2 \vee s_{2,1}$	$s_{1,1} \lor \overline{s}_{2,2}$	$\overline{s}_{2,2} \lor s_{3,2}$	$s_{2,2} \vee \overline{s}_{3,3}$	
$\overline{s}_{1,1} \lor s_{2,1}$	$\overline{\ell}_3 \lor s_{3,1}$	$\ell_3 \vee s_{2,2} \vee \overline{s}_{3,2}$	$s_{3,2}$	
$\ell_2 \vee s_{1,1} \vee \overline{s}_{2,1}$	$\overline{s}_{2,1} \lor s_{3,1}$	$s_{2,1} \vee \overline{s}_{3,2}$		

![](_page_15_Figure_1.jpeg)

l	$_1+\ell_2+\ell_3\geq 2$	<i>Si</i> -1, <i>j</i>		Block 1 Block 2	Block 3
ℕ <i>s</i> ,	leaning of $s_{i,j} ext{-variable:}$ $s_{i,j}  ext{ true} \Leftrightarrow \ell_1 + \ldots + \ell_i \geq j$	$s_{i-1,j-1}$			$l_3$
ſ	$\overline{\ell}_1 \lor s_{1,1}$ $\ell_1 \lor \overline{s}_{1,1}$	$\overline{\ell}_2 \vee \overline{s}_{1,1} \vee s_{2,2}$ $\ell_2 \vee \overline{s}_{2,2}$	$\ell_3 \lor s_{2,1} \lor \overline{s}_{3,1}$ $\overline{\ell}_2 \lor \overline{s}_{2,1} \lor s_{2,2}$	$\overline{\ell}_3 \lor \overline{s}_{2,2} \lor s_{3,3}$ $\ell_2 \lor \overline{s}_{2,2}$	
	$\overline{\ell}_2 \lor s_{2,1}$	$s_2 \vee s_{2,2}$ $s_{1,1} \vee \overline{s}_{2,2}$ $\overline{\ell}_2 \vee s_2$	$\overline{s}_{2,2} \lor \overline{s}_{3,2}$	$s_{2,2} \lor \overline{s}_{3,3}$	
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## Pseudo-Boolean Proof Logging with $\operatorname{Ver}\operatorname{IPB}$

- Inspired by SAT proof logging, but operates on 0--1 linear constraints
- Supports efficient proof logging for
  - ► SAT solving including advanced techniques previously beyond efficient proof logging like
    - ★ Gaussian elimination [GN21]
    - ★ Symmetry breaking [BGMN23]
  - SAT-based optimization (MaxSAT) approaches like
    - ★ Model-improving search [VDB22]
    - ★ Core-guided search [BBN<sup>+</sup>23]
  - Subgraph problems [GMN20, GMM<sup>+</sup>20]
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This work:

- Proof logging for translating pseudo-Boolean constraints to CNF
- General framework to certify many different encodings

## End-to-End Verification for Pseudo-Boolean Solving

![](_page_19_Figure_1.jpeg)

### All the Technical Details on One Slide

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How pretty much all PB-to-CNF translations work:

- **(**) Design circuit evaluating left-hand side of 0-1 integer linear constraint
- **2** Encode circuit to CNF using Tseitin translation (one new variable per circuit wire)
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We solve this task by

- Understanding the circuit design
- Oeriving CNF encoding of circuit (easy)
- Proving circuit output is true (tricky, but using pseudo-Boolean reasoning really helps)

## **Experimental Evaluation**

- $\bullet$  Certified translations for CNF encodings with  $\mathrm{VerlTAsPBLiB}^1$ 
  - Sequential counter [Sin05]
  - Totalizer [BB03]
  - Generalized totalizer [JMM15]
  - Adder network [ES06]
- $\bullet$  Proofs verified by proof checker  $\rm VERIPB^2$
- $\bullet$  Formulas solved with fork of  ${\rm KISSAT}^3$  outputting  ${\rm VERIPB}$  proofs
- Benchmarks from PB 2016 Evaluation<sup>4</sup> in 3 categories
  - Only cardinality constraints (sequential counter, totalizer)
  - ► Only general 0-1 ILP constraints (generalized totalizer, adder network)
  - ► Mixed cardinality & general 0-1 ILP constraints (sequential counter + adder network)

<sup>1</sup>https://github.com/forge-lab/VeritasPBLib
<sup>2</sup>https://gitlab.com/MIAOresearch/software/VeriPB
<sup>3</sup>https://gitlab.com/MIAOresearch/tools-and-utilities/kissat\_fork
<sup>4</sup>http://www.cril.univ-artois.fr/PB16/

CNF Size vs Proof Size in KiB

![](_page_25_Figure_1.jpeg)

• Nice scaling for proof size in terms of original CNF formula size

• Except for some sequential encoding cases (which is not such a great encoding anyway)

## Translation Time vs Proof Checking Time in Seconds

![](_page_26_Figure_1.jpeg)

• Proof checking slower — has to verify full proof

Jakob Nordström (UCPH & LU)

# Solving Time vs Proof Checking Time in Seconds

![](_page_27_Figure_1.jpeg)

- Room for improvement of end-to-end proof checking process
- But even first proof-of-concept implementation shows our approach is viable

Jakob Nordström (UCPH & LU)

### Future Work

Improve performance and reliability:

- Cutting planes derivations instead of reverse unit propagations [VDB22]
- Backwards checking/trimming for verification (as in DRAT-TRIM [HHW13a])
- Fully formally verified proof checking (work in progress [BMM<sup>+</sup>23])

### Extend proof logging further:

- PB-to-CNF translations with odd-even mergesort & bitonic sorting networks [Bat68]
- All of MaxSAT solving and (linear) pseudo-Boolean optimization
- Mixed integer linear programming
- Automated planning

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We're hiring! Talk to me to join the proof logging revolution! ©

This work:

- General approach for certifying many different PB-to-CNF translations
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- Boolean satisfiability (SAT) including advanced techniques [GN21, BGMN23]
- SAT-based optimization (MaxSAT) [VDB22, BBN<sup>+</sup>23]
- SAT-based pseudo-Boolean solving [this work]
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Action point: What can  $\mathrm{Ver}\mathrm{iPB}$  do for you?  $\odot$ 

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## Thank you for your attention!

Jakob Nordström (UCPH & LU)

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